
On avoiding dependent choices in Formal Topology

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.

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Why avoid countable/dependent choice?

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- *“The theory should be **compatible** with known theories such as classical set theory and the theory of a generic topos, ... and hence it should be **minimal** with respect to such more expressive existing theories”* - from “Toward a minimalist foundation for constructive mathematics”, by Maria Emilia Maietti and Giovanni Sambin.

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- I suggest: A good **core (minimal?)** theory for constructive set theory should, at least, be preserved by the Heyting valued model construction.

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For 1,2 $uREA + DC$ can be weakened to $*REA$.

For 3 something that seems a bit stronger than $*REA$, $*_2REA$, seems to be needed.

Plan of talk

- Review
- Relation Reflection
- Coinductive Definitions
- If time: a slide on characterising the collection of positivity relations

Some axiom systems for Constructive Set Theory

These are formulated in the language of ZF, but use intuitionistic logic.

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- $CZF^+ \equiv CZF^o +$ Regular Extension Axiom.

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- **Restricted Separation:** For each restricted class X , if a is a set the class $a \cap X$ is a set.
- A class is **restricted** if it is defined by a restricted formula; i.e. a formula all of whose quantifiers have one of the forms $(\forall x \in t)$ or $(\exists x \in t)$ where t is a variable or parameter.

The Collection Principles

- **Definition:** For classes X, Y, R

$$R : X \succ Y \quad \text{iff} \quad (\forall x \in X)(\exists y \in Y) (x, y) \in R$$

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- **or:** For all sets a, b there is a set c of subsets of $a \times b$ such that for every set $r : a \succ b$ there is a set $r' \in c$ such that $r \supseteq r' : a \succ b$.

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Countable Choice (CC): For all sets X, R , if $R : \mathbb{N} \multimap X$ then there is a $f : \mathbb{N} \rightarrow X$ such that

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Inductive Definitions

- Let C be a **covering system** on a class S ; i.e. an operation $C : S \rightarrow Pow(Pow(S))$.
- X/a is a **C -step** if $a \in S$ and $X \in C(a)$.
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Theorem ($CZF^0 + REA$): If S is a set, for each set $U \subseteq S$ the class $\mathcal{A}U$ is a set. Moreover there is a covering system D on S such that such that, for all $U \in Pow(S)$, if $a \in S$ then

$$a \in \mathcal{A}U \Leftrightarrow (\exists V \in D(a))[V \subseteq U].$$

Coinductive Definitions

Let C be a covering system on a class S . A class $U \subseteq S$ is **C -progressive** if, for all C -steps X/a ,

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Theorem ($CZF^o + RDC$): For each class $U \subseteq S$ there is a largest C -progressive class $\mathcal{J}U$ included in U . $\mathcal{J}U$ is the class **coinductively defined by C, U** .

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Theorem ($CZF^o + uREA + DC$): If S is a set then, for each set $U \subseteq S$, the class $\mathcal{J}U$ exists and is a set.

● $uREA + DC$ can be replaced by $*REA$.

An example avoiding Dependent Choices, 1

Theorem [$CZF^o + ?$]: Let C be a covering system on a class S and let

$$J = \bigcup \{V \in Pow(S) \mid V \text{ is } C\text{-progressive}\}.$$

Then J is the largest C -progressive class; i.e. (i) J is C -progressive, (ii) If B is a C -progressive class then $B \subseteq J$.

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Proof: (i) is easy. For (ii), for classes B_1, B_2 let

$$B_1 \mapsto B_2 \equiv \forall a \in B_1 \forall X \in C(a) X \cap B_2.$$

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So $\forall a \in B \forall X \in C(a) \exists y \in B y \in X$,

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So $\forall a \in B \forall X \in C(a) \exists y \in B y \in X$, and so, by Strong Collection,

$$(*) \quad \forall a \in B \exists Y \in Pow(B) \forall X \in C(a) \exists y \in Y y \in X.$$

An example avoiding Dependent Choices, 1

Theorem [$CZF^o+?$]: Let C be a covering system on a class S and let

$$J = \bigcup \{V \in Pow(S) \mid V \text{ is } C\text{-progressive}\}.$$

Then J is the largest C -progressive class; i.e. (i) J is C -progressive, (ii) If B is a C -progressive class then $B \subseteq J$.

Proof: (i) is easy. For (ii), for classes B_1, B_2 let

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We must show that $\forall a \in B \exists V \in Pow(S) [a \in V \mapsto V]$.

An example avoiding Dependent Choices, 2

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To prove: $\forall a \in B \exists V \in Pow(S)[a \in V \mapsto V]$. Let $a \in B$.

Then $\{a\} \in Pow(B)$ so that, by the lemma and **RDC** there is $f : \mathbb{N} \rightarrow Pow(B)$ such that $f(0) = \{a\}$ and, for $n \in \mathbb{N}$,

$$f(n) \mapsto f(n + 1).$$

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Let $V = \bigcup_{n \in \mathbb{N}} f(n) \in Pow(S)$. Then $a \in V$. Also $V \mapsto V$ as

$$\begin{aligned} b \in V &\Rightarrow b \in f(n) \text{ for some } n \in \mathbb{N} \\ &\Rightarrow \forall X \in C(b) X \not\mapsto f(n+1) \\ &\Rightarrow \forall X \in C(b) X \not\mapsto V, \end{aligned}$$

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Using **RRS** instead of **RDC**: there is a set $\mathcal{X} \subseteq Pow(B)$ such that $\{a\} \in \mathcal{X}$ and $\forall U \in \mathcal{X} \exists V \in \mathcal{X} U \mapsto V$. Let

$V = \bigcup \mathcal{X} \in Pow(S)$. Then $a \in V \mapsto V$.

An example avoiding Dependent Choices, 3

Proof of Lemma: $\forall U \in Pow(B) \exists V \in Pow(B) U \mapsto V.$

Recall that

$$(*) \quad \forall a \in B \exists Y \in Pow(B) \forall X \in C(a) Y \not\supseteq X.$$

An example avoiding Dependent Choices, 3

Proof of Lemma: $\forall U \in Pow(B) \exists V \in Pow(B) U \mapsto V.$

Recall that

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Proof of Lemma: $\forall U \in Pow(B) \exists V \in Pow(B) U \mapsto V.$

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Let $U \in Pow(B)$. Then

$$\forall a \in U \exists Y \in Pow(B) \forall X \in C(a) Y \not\subseteq X.$$

By Strong Collection there is a set $\mathcal{Z} \subseteq Pow(B)$ such that

$$\forall a \in U \exists Y \in \mathcal{Z} \forall X \in C(a) Y \not\subseteq X.$$

Let $V = \cup \mathcal{Z}$. Then $U \mapsto V$.

Conclusion

Recall:

$$CZF^+ \equiv CZF^o + REA$$

$$CZF^u \equiv CZF^o + uREA$$

$$CZF^* \equiv CZF^o + *REA$$

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- CZF^* and CZF^{*2} can be used to prove certain results that had been previously proved in $CZF^u + DC$.
- **Conjecture:** CZF^* (and CZF^{*2}) are preserved in Heyting valued models over set-presented cHa's.
- **But** DC does not generally hold in Heyting valued models.

A characterisation of the positivity relations \bowtie on (S, \triangleleft)

Let $\mathcal{S} = (S, \triangleleft)$ be a fixed formal topology. I prefer to work with the **positivity operators** $\mathcal{J} : Pow(S) \rightarrow Pow(S)$ where, for all a, U , $a \in \mathcal{J}U \equiv a \bowtie U$.

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- Let I be a class of completely prime sets. For all U let

$$\mathcal{J}_I U = \cup \{F \in I \mid F \subseteq U\},$$

and call I **standard** if

$(\forall I' \in Pow(I)) \cup I' \in I$ and $(\forall U \in Pow(S)) \mathcal{J}_I U$ is a set.

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Note: \mathcal{J}_I is a **set-presented** positivity operator for any **set** I of completely prime sets.

x0

x1

 x2

x0

x1

● x2

● x3

x0

x1

 x2

 x3

x4

x0

x1

 x2

 x3

x4